On the solutions of a max-type system of

difference equations of higher order

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Abstract

In this paper, we study the following max-type system of difference equations of higher order:

$$\begin{cases} x_n = \max\{A, \frac{y_{n-t}}{x_{n-s}}\}, \\ y_n = \max\{B, \frac{x_{n-t}}{y_{n-s}}\}, \end{cases} \quad n \in \{0, 1, 2, \dots\}, \end{cases}$$

where $A, B \in (0, +\infty)$, $t, s \in \{1, 2, ...\}$ with gcd(s, t) = 1, the initial values $x_{-d}, y_{-d}, x_{-d+1}, y_{-d+1}, ..., x_{-1}, y_{-1} \in (0, +\infty)$ and $d = max\{t, s\}$.

MSC: 39A10; 39A11

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1 Introduction

Concrete nonlinear difference equations and systems have attracted some recent attention (see, e.g., [1–39]). One of the classes of such equations/systems are max-type difference equations/systems. For some results of solutions of many max-type difference equations and systems, such as eventual periodicity, the boundedness character and attractivity, see, e.g. [1–5, 7–9, 11–16, 18–25, 28–30, 32–36, 38, 39] and the references therein. Our purpose in this paper is to study the eventual periodicity of the following max-type system of difference equation of higher order:

$$\begin{cases} x_n = \max\{A, \frac{y_{n-t}}{x_{n-s}}\}, \\ y_n = \max\{B, \frac{x_{n-t}}{y_{n-s}}\}, \end{cases} \quad n \in \mathbf{N}_0 \equiv \{0, 1, \ldots\},$$
(1.1)

where $A, B \in \mathbf{R}_+ \equiv (0, +\infty), t, s \in \mathbf{N} \equiv \{1, 2, ...\}$ with gcd(s, t) = 1, the initial values $x_{-d}, y_{-d}, x_{-d+1}, y_{-d+1}, ..., x_{-1}, y_{-1} \in \mathbf{R}_+$ and $d = \max\{t, s\}$.

When t = 1 and s = 2, (1.1) reduces to the max-type system of difference equations

$$\begin{cases} x_n = \max\{A, \frac{y_{n-1}}{x_{n-2}}\}, \\ y_n = \max\{B, \frac{x_{n-1}}{y_{n-2}}\}, \end{cases} & n \in \mathbf{N}_0. \end{cases}$$
(1.2)

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Fotiades and Papaschinopoulos in [5] showed that every positive solution of (1.2) is eventually periodic.

In 2012, Stević [23] obtained in an elegant way the general solution to the following max-type system of difference equations:

$$\begin{cases} x_{n+1} = \max\{\frac{A}{x_n}, \frac{y_n}{x_n}\}, \\ y_{n+1} = \max\{\frac{A}{y_n}, \frac{x_n}{y_n}\}, \end{cases} \quad n \in \mathbf{N}_0, \end{cases}$$
(1.3)

for the case $x_0, y_0 \ge A > 0$ and $y_0/x_0 \ge \max\{A, 1/A\}$.

In [35], Sun and Xi studied the following max-type system of difference equations:

$$\begin{cases} x_n = \max\{\frac{1}{x_{n-m}}, \min\{1, \frac{A}{y_{n-r}}\}\},\\ y_n = \max\{\frac{1}{y_{n-m}}, \min\{1, \frac{B}{x_{n-t}}\}\}, \end{cases} \quad n \in \mathbf{N}_0, \end{cases}$$
(1.4)

where $A, B \in \mathbf{R}_+$, $m, r, t \in \mathbf{N}$ and the initial values $x_{-d}, y_{-d}, x_{-d+1}, y_{-d+1}, \dots, x_{-1}, y_{-1} \in \mathbf{R}_+$ with $d = \max\{m, r, t\}$ and showed that every positive solution of (1.4) is eventually periodic with period 2m.

When m = r = t = 1 and A = B, (1.4) reduces to the max-type system of difference equations

$$\begin{cases} x_n = \max\{\frac{1}{x_{n-1}}, \min\{1, \frac{A}{y_{n-1}}\}\},\\ y_n = \max\{\frac{1}{y_{n-1}}, \min\{1, \frac{A}{x_{n-1}}\}\}, \end{cases} \quad n \in \mathbf{N}_0. \end{cases}$$
(1.5)

Yazlik et al. [39] in 2015 obtained in an elegant way the general solution of (1.5).

In 2012, Stević [24] studied the following max-type system of difference equations:

$$\begin{cases} y_{n}^{(1)} = \max_{1 \le i \le m_{1}} \{ f_{1i}(y_{n-k_{i,1}^{(1)}}^{(1)}, y_{n-k_{i,2}^{(2)}}^{(2)}, \dots, y_{n-k_{i,l}^{(1)}}^{(l)}, n), y_{n-s}^{(1)} \}, \\ y_{n}^{(2)} = \max_{1 \le i \le m_{2}} \{ f_{2i}(y_{n-k_{i,1}^{(2)}}^{(1)}, y_{n-k_{i,2}^{(2)}}^{(2)}, \dots, y_{n-k_{i,l}^{(2)}}^{(l)}, n), y_{n-s}^{(2)} \}, \\ \dots, \\ y_{n}^{(l)} = \max_{1 \le i \le m_{l}} \{ f_{li}(y_{n-k_{i,1}^{(l)}}^{(1)}, y_{n-k_{i,2}^{(l)}}^{(2)}, \dots, y_{n-k_{i,l}^{(l)}}^{(l)}, n), y_{n-s}^{(l)} \}, \end{cases}$$
(1.6)

where $s, l, m_j, k_{i,t}^{(j)} \in \mathbf{N}$ $(j, t \in \{1, 2, ..., l\})$ and $f_{ji} : \mathbf{R}_+^l \times \mathbf{N}_0 \longrightarrow \mathbf{R}_+$ $(j \in \{1, ..., l\}$ and $i \in \{1, ..., m_j\}$, and showed that every positive solution of (1.6) is eventually periodic with (not necessarily prime) period *s* if f_{ji} satisfy some conditions.

Moreover, Stević et al. [29] in 2014 investigated the following max-type system of difference equations:

$$\begin{cases} y_n^{(1)} = \max_{1 \le i_1 \le m_1} \{ f_{1i_1}(y_{n-k_{i_{1,1}}^{(1)}}, y_{n-k_{i_{1,2}}^{(2)}}^{(2)}, \dots, y_{n-k_{i_{1,l}}^{(1)}}^{(l)}, n), y_{n-t_1s}^{(\sigma(1))} \}, \\ x_n^{(2)} = \max_{1 \le i_2 \le m_2} \{ f_{2i_2}(y_{n-k_{i_{2,2}}^{(2)}}, y_{n-k_{i_{2,2}}^{(2)}}^{(2)}, \dots, y_{n-k_{i_{2,l}}^{(2)}}^{(l)}, n), y_{n-t_{2s}}^{(\sigma(2))} \}, \\ \dots, \\ y_n^{(l)} = \max_{1 \le i_l \le m_l} \{ f_{li_l}(y_{n-k_{i_{l,1}}^{(l)}}, y_{n-k_{i_{l,2}}^{(l)}}^{(2)}, \dots, y_{n-k_{i_{l,l}}^{(l)}}^{(l)}, n), y_{n-t_{ls}}^{(\sigma(l))} \}, \end{cases}$$
(1.7)

where $s, l, m_j, t_j, k_{i_j,h}^{(j)} \in \mathbf{N}$ $(j, h \in \{1, 2, ..., l\}), (\sigma(1), ..., \sigma(l))$ is a permutation of (1, ..., l) and $f_{ji_j} : \mathbf{R}_+^l \times \mathbf{N}_0 \longrightarrow \mathbf{R}_+$ $(j \in \{1, ..., l\}$ and $i_j \in \{1, ..., m_j\})$. They showed that every positive solution of (1.7) is eventually periodic with period sT for some $T \in \mathbf{N}$ if f_{ji_j} satisfy some conditions.

2 Main results and proofs

In this section, we study the eventual periodicity of positive solutions of system (1.1). Let $\{(x_n, y_n)\}_{n \ge -d}$ be a solution of (1.1) with the initial values $x_{-d}, y_{-d}, x_{-d+1}, y_{-d+1}, \dots, x_{-1}, y_{-1} \in \mathbf{R}_+$.

Lemma 2.1 If $x_n = A$ eventually, then y_n is a periodic sequence with period 2s eventually. If $y_n = B$ eventually, then x_n is a periodic sequence with period 2s eventually.

Proof Assume that $x_n = A$ eventually. By (1.1) we see

$$y_n = \max\left\{B, \frac{A}{y_{n-s}}\right\}$$
 eventually, (2.1)

which implies $y_n y_{n-s} \ge A$ eventually and

$$B \leq y_n = \max\left\{B, \frac{A}{y_{n-s}}\right\}$$
$$= \max\left\{B, \frac{Ay_{n-2s}}{y_{n-s}y_{n-2s}}\right\}$$
$$\leq \max\{B, y_{n-2s}\} \leq y_{n-2s} \quad \text{eventually.}$$
(2.2)

Then, for any $0 \le i \le 2s - 1$, y_{2ns+i} is eventually nonincreasing.

We claim that, for every $0 \le i \le 2s - 1$, y_{2ns+i} is a constant sequence eventually. Assume on the contrary that, for some $0 \le i \le 2s - 1$, y_{2ns+i} is not a constant sequence eventually. Then there exists a sequence of positive integers $k_1 < k_2 < \cdots$ such that, for any $n \in \mathbb{N}$, we have

$$B < y_{2sk_{n+1}+i} = \frac{A}{y_{2sk_{n+1}+i-s}}$$

$$< y_{2sk_n+i} = \frac{A}{y_{2sk_n+i-s}},$$
 (2.3)

which implies $y_{2sk_{n+1}+i-s} > y_{2sk_n+i-s}$ for any $n \in \mathbb{N}$. This is a contradiction. Thus y_n is a periodic sequence with period 2*s* eventually. The second case follows from the previously proved one by interchanging letters. The proof is complete.

Lemma 2.2 If $A \ge B \ge 1/A$, then $x_{2(n+1)t+i} \le x_{2nt+i}$ for any $n \ge t + s$ and $i \in \mathbb{N}_0$. If $B \ge A \ge 1/B$, then $y_{2(n+1)t+i} \le y_{2nt+i}$ for any $n \ge t + s$ and $i \in \mathbb{N}_0$.

Proof Assume that $A \ge B \ge 1/A$. By (1.1) we see that $x_n \ge A$ and $y_n \ge B$ for any $n \in \mathbb{N}_0$, and

$$x_{2(n+1)t+i} = \max\left\{A, \frac{B}{x_{2(n+1)t+i-s}}, \frac{x_{2nt+i}}{x_{2(n+1)t+i-s}y_{2(n+1)t+i-t-s}}\right\}.$$
(2.4)

Since $B/x_{2(n+1)t+i-s} \le B/A \le 1 \le A$ and $x_{2(n+1)t+i-s}y_{2(n+1)t+i-t-s} \ge AB \ge 1$ for $2(n+1)t+i \ge t+s$, we obtain

$$x_{2(n+1)t+i} \le \max\{A, x_{2nt+i}\} \le x_{2nt+i}.$$
(2.5)

The second case follows from the previously proved one by interchanging letters. The proof is complete. $\hfill \Box$

Theorem 2.1 Let AB > 1. If $A \ge B$, then $x_n = A$ eventually and y_n is a periodic sequence with period 2s eventually. If B > A, then $y_n = B$ eventually and x_n is a periodic sequence with period 2s eventually.

Proof Assume that $A \ge B$. For any $0 \le i \le 2t - 1$ and $n \in \mathbb{N}_0$, we have

$$A \le x_{2(n+1)t+i} = \max\left\{A, \frac{x_{2nt+i}}{x_{2(n+1)t-s+i}y_{2(n+1)t-s-t+i}}\right\}.$$
(2.6)

By Lemma 2.2 we may let $\lim_{n\to\infty} x_{2nt+i} = A_i$. Note that

$$\lim_{n \to \infty} \frac{x_{2nt+i}}{x_{2(n+1)t-s+i}y_{2(n+1)t-s-t+i}} \le \lim_{n \to \infty} \frac{x_{2nt+i}}{AB} = \frac{A_i}{AB} < A_i$$

$$(2.7)$$

and

$$\lim_{n \to \infty} x_{2(n+1)t+i} = A_i. \tag{2.8}$$

Thus we have $x_n = A$ eventually. By Lemma 2.1, we see that y_n is a periodic sequence with period 2*s* eventually. The second case follows from the previously proved one by interchanging letters. The proof is complete.

In the following, we assume AB = 1. For any $i \in \mathbb{N}_0$, let

$$\lim_{n \to \infty} x_{2nt+i} = A_i \quad \text{if } A \ge B \tag{2.9}$$

and

$$\lim_{n \to \infty} y_{2nt+i} = B_i \quad \text{if } B \ge A. \tag{2.10}$$

Then $A_i \ge A$ and $B_i \ge B$.

Lemma 2.3 If $A \ge B = 1/A$ and $A_i > A$ for some $i \in \mathbb{N}_0$, then, for any $k \in \mathbb{N}$, $x_{2nt+ks+i}$ and $y_{2nt-t+ks+i}$ are constant sequences eventually. If $B \ge A = 1/B$ and $B_i > B$ for some $i \in \mathbb{N}_0$, then, for any $k \in \mathbb{N}$, $y_{2nt+ks+i}$ and $x_{2nt-t+ks+i}$ are constant sequences eventually.

Proof Assume that $A \ge B$ and $A_i > A$ for some $i \in \mathbb{N}_0$. Since $A_i > A$, it follows from Lemma 2.2 and (1.1) that

$$x_{2nt+i} = \max\left\{A, \frac{x_{2(n-1)t+i}}{x_{2nt-s+i}y_{2nt-t-s+i}}\right\}$$

= $\frac{x_{2(n-1)t+i}}{x_{2nt-s+i}y_{2nt-t-s+i}}$ eventually. (2.11)

From this we have

$$B \leq \lim_{n \to \infty} y_{2nt-t-s+i}$$

$$= \lim_{n \to \infty} \frac{x_{2(n-1)t+i}}{x_{2nt+i}x_{2nt-s+i}}$$

$$= \frac{1}{A_{-s+i}} \leq \frac{1}{A} = B.$$
(2.12)

This implies

$$\lim_{n \to \infty} x_{2nt-s+i} = A \tag{2.13}$$

and

$$\lim_{n \to \infty} y_{2nt-t-s+i} = B \tag{2.14}$$

and

$$\lim_{n \to \infty} y_{2nt-t+i} = \lim_{n \to \infty} x_{2nt+i} x_{2nt-s+i} = A_i A.$$
(2.15)

Since

$$x_{2nt+s+i} = \max\left\{A, \frac{x_{2(n-1)t+s+i}}{x_{2nt+i}y_{2nt-t+i}}\right\}$$
(2.16)

and

$$\lim_{n \to \infty} \frac{x_{2(n-1)t+s+i}}{x_{2nt+i}y_{2nt-t+i}} = \frac{A_{s+i}}{A_i^2 A} < A_{s+i},$$
(2.17)

we see that $x_{2nt+s+i} = A$ eventually. Note that

$$\lim_{n \to \infty} \frac{x_{2(n-1)t+s+i}}{y_{2nt-t+i}} = \frac{A}{AA_i} < B,$$
(2.18)

from which it follows that

$$y_{2nt-t+s+i} = \max\left\{B, \frac{x_{2(n-1)t+s+i}}{y_{2nt-t+i}}\right\}$$
$$= B \quad \text{eventually}, \tag{2.19}$$

and

$$y_{2nt-t+2s+i} = \max\left\{B, \frac{x_{2(n-1)t+2s+i}}{y_{2nt-t+s+i}}\right\}$$
$$= \frac{x_{2(n-1)t+2s+i}}{B} \quad \text{eventually,}$$
(2.20)

and

$$x_{2nt+2s+i} = \max\left\{A, \frac{y_{2nt-t+2s+i}}{x_{2nt+s+i}}\right\}$$

= max{A, x_{2(n-1)t+2s+i}}
= x_{2(n-1)t+2s+i} eventually. (2.21)

If $x_{2nt+2s+i} > A$ eventually, then, in a similar fashion, we obtain:

(1) $x_{2nt+3s+i} = A$ eventually and $y_{2nt-t+3s+i} = B$ eventually.

(2) $x_{2nt+4s+i}$ and $y_{2nt-t+4s+i}$ are constant sequences eventually.

If $x_{2nt+2s+i} = A$ eventually, then $y_{2nt-t+2s+i} = A/B$ eventually, and

$$y_{2nt-t+3s+i} = \max\left\{B, \frac{x_{2(n-1)t+3s+i}}{y_{2nt-t+2s+i}}\right\}$$

= $\max\left\{B, \frac{x_{2(n-1)t+3s+i}B}{A}\right\}$
= $\frac{x_{2(n-1)t+3s+i}B}{A}$ eventually, (2.22)

and

$$x_{2nt+3s+i} = \max\left\{A, \frac{y_{2nt-t+3s+i}}{x_{2nt+2s+i}}\right\}$$
$$= \max\left\{A, \frac{x_{2(n-1)t+3s+i}B}{A^2}\right\} \quad \text{eventually.}$$
(2.23)

From this we see that if A = B, then

 $x_{2nt+3s+i} = x_{2(n-1)t+3s+i}$ eventually, (2.24)

and if A > B, then

$$x_{2nt+3s+i} = A \quad \text{eventually,} \tag{2.25}$$

since

$$\lim_{n \to \infty} \frac{x_{2(n-1)t+3s+i}B}{A^2} = \frac{A_{3s+i}B}{A^2} < A_{3s+i}.$$
(2.26)

Using induction and arguments similar to the ones developed in the above given proof, we can show that, for any $k \in \mathbf{N}$, $x_{2nt+ks+i}$ and $y_{2nt-t+ks+i}$ are constant sequences eventually. The second case follows from the previously proved one by interchanging letters. The proof is complete.

Lemma 2.4 If A = 1/B > B and for some $i \in \mathbb{N}_0$, $x_{2nt+i} > A$ eventually and $A_i = A$, then, for any $k \in \mathbb{N}$, $x_{2nt+ks+i}$ and $y_{2nt-t+ks+i}$ are constant sequences eventually. If B = 1/A > A and for some $i \in \mathbb{N}_0$, $y_{2nt+i} > B$ eventually and $B_i = B$, then, for any $k \in \mathbb{N}$, $y_{2nt+ks+i}$ and $x_{2nt-t+ks+i}$ are constant sequences eventually.

Proof Assume that A = 1/B > B and for some $i \in \mathbb{N}_0$, $x_{2nt+i} > A$ eventually and $A_i = A$. By (1.1) we have

$$x_{2nt+i} = \max\left\{A, \frac{y_{2nt-t+i}}{x_{2nt-s+i}}\right\}$$
$$= \frac{y_{2nt-t+i}}{x_{2nt-s+i}} \quad \text{eventually,}$$
(2.27)

and

$$x_{2nt+s+i} = \max\left\{A, \frac{x_{2(n-1)t+s+i}}{x_{2nt+i}y_{2nt-t+i}}\right\}$$
(2.28)

and

$$\lim_{n \to \infty} \frac{x_{2(n-1)t+s+i}}{x_{2nt+i}y_{2nt-t+i}} = \lim_{n \to \infty} \frac{x_{2(n-1)t+s+i}}{x_{2nt+i}^2 x_{2nt-s+i}}$$
$$\leq \frac{A_{s+i}}{A^3} < A_{s+i}.$$
(2.29)

Then we see that $x_{2nt+s+i} = A$ eventually. From this and $y_{2nt-t+i} \ge A^2$ eventually it follows that

$$y_{2nt-t+s+i} = \max\left\{B, \frac{x_{2(n-1)t+s+i}}{y_{2nt-t+i}}\right\}$$
$$= \max\left\{B, \frac{A}{y_{2nt-t+i}}\right\}$$
$$= B \quad \text{eventually}, \tag{2.30}$$

and

$$y_{2nt-t+2s+i} = \max\left\{B, \frac{x_{2(n-1)t+2s+i}}{y_{2nt-t+s+i}}\right\} = \frac{x_{2(n-1)t+2s+i}}{B} \quad \text{eventually,}$$
(2.31)

and

$$x_{2nt+2s+i} = \max\left\{A, \frac{y_{2nt-t+2s+i}}{x_{2nt+s+i}}\right\}$$

= max{A, x_{2(n-1)t+2s+i}}
= x_{2(n-1)t+2s+i} eventually. (2.32)

Thus $x_{2nt+2s+i}$ and $y_{2nt-t+2s+i}$ are constant sequences eventually. Using arguments similar to the ones developed in the proof of Lemma 2.3, we can show that, for any $k \in \mathbf{N}$, $x_{2nt+ks+i}$ and

 $y_{2nt-t+ks+i}$ are constant sequences eventually. The second case follows from the previously proved one by interchanging letters. The proof is complete.

Theorem 2.2

- (1) Assume A = 1/B > B. Then one of the following statements holds.
 - (i) $x_n = A$ eventually and y_n is a periodic sequence with period 2s eventually.
 - (ii) If s is odd, then x_n , y_n are periodic sequences with period 2t eventually.
 - (iii) If s is even, then x_n is a periodic sequence with period 2t eventually and y_n is a periodic sequence with period 2st eventually.
- (2) Assume B = 1/A > A. Then one of the following statements holds.
 - (i) $y_n = B$ eventually and x_n is a periodic sequence with period 2s eventually.
 - (ii) If s is odd, then x_n , y_n are periodic sequences with period 2t eventually.
 - (iii) If s is even, then y_n is a periodic sequence with period 2t eventually and x_n is a periodic sequence with period 2st eventually.

Proof Assume that A = 1/B > B. If $x_n = A$ eventually, then by Lemma 2.1 we see that y_n is a periodic sequence with period 2*s* eventually. Now we assume that $x_n \neq A$ eventually. Then we have $A_i > A$ (or $x_{2nt+i} > A$ eventually and $A_i = A$) for some $0 \le i \le 2t - 1$.

If *s* is odd, then gcd(2t, s) = 1. Thus, for every $j \in \{0, 1, 2, ..., 2t - 1\}$, there exist some $1 \le i_j \le 2t$ and integer λ_j such that $i_j s = \lambda_j 2t + j$ since $\{rs : 0 \le r \le 2t - 1\} = \{0, 1, 2, ..., 2t - 1\}$ (mod 2*t*). By Lemma 2.3 and Lemma 2.4 we see that, for any $k \in \mathbb{N}$, $x_{2nt+ks+i}$ and $y_{2nt-t+ks+i}$ are constant sequences eventually. Thus, for any $0 \le r \le 2t - 1$, x_{2nt+r} and y_{2nt+r} are constant sequences eventually, which implies that x_n , y_n are periodic sequences with period 2*t* eventually.

In the following, we assume that *s* is even with s = 2s'. Then gcd(t,s') = 1 and *t* is odd. Thus, for every $j \in \{0, 1, 2, ..., t - 1\}$, there exist some $1 \le i_j \le t$ and integer λ_j such that $i_js' = \lambda_jt + j$ and $i_js = \lambda_j2t + 2j$.

If $x_{2nt+i} \neq A$ eventually for some $i \in \{0, 2, ...\}$ and $x_{2nt+l} \neq A$ eventually for some $l \in \{1, 3, ...\}$, then by Lemma 2.3 and Lemma 2.4 we see that, for any $k \in \mathbb{N}$, $x_{2nt+ks+i}$, $y_{2nt-t+ks+i}$, $x_{2nt+ks+l}$ and $y_{2nt-t+ks+l}$ are constant sequences eventually. Thus, for any $0 \leq r \leq 2t - 1$, x_{2nt+r} and y_{2nt+r} are constant sequences eventually, which implies that x_n , y_n are periodic sequences with period 2t eventually.

If $x_{2nt+i} \neq A$ eventually for some $i \in \{0, 2, ...\}$ and $x_{2nt+l} = A$ eventually for any $l \in \{1, 3, ...\}$, then by Lemma 2.3 and Lemma 2.4 we see that, for any $k \in \mathbb{N}$, $x_{2nt+ks+i}$ and $y_{2nt-t+ks+i}$ are constant sequences eventually. This implies that, for every $r \in \{0, 1, 2, ..., 2t - 1\}$, x_{2nt+r} is constant sequence eventually and for every $l \in \{1, 3, ...\}$, y_{2nst+l} is constant sequence eventually. By (1.1) we see that there exists $N \in \mathbb{N}$ such that, for any $n \ge N$ and $r \in \{0, 2, ...\}$,

$$y_{2nt+r} = \max\left\{B, \frac{A}{y_{2nt+r-s}}\right\}.$$
 (2.33)

Then we have $y_{2nt+r}y_{2nt+r-s} \ge A$. Thus, for any $n \ge N$ and $l \in \{1, 3, ...\}$ and $k \in \mathbb{N}$,

$$B \le y_{2nt+t+2ks+l} = \max\left\{B, \frac{x_{2nt+2ks+l}}{y_{2nt+t+2ks+l-s}}\right\}$$
$$= \max\left\{B, \frac{Ay_{2nt+t+2ks+l-2s}}{y_{2nt+t+2ks+l-s}y_{2nt+t+2ks+l-2s}}\right\}$$

$$\leq \max\{B, y_{2nt+t+2ks-2s+l}\}$$

$$= y_{2nt+t+2ks-2s+l} \quad \text{eventually.} \tag{2.34}$$

Then, for every $n \ge N$ and $l \in \{1, 3, \ldots\}$, we have

$$B \le \dots \le y_{2nt+t+2ks+l} \le y_{2nt+t+2ks-2s+l} \le y_{2nt+t+2s+l} \le y_{2nt+t+l}.$$
(2.35)

We claim that, for every $n \ge N$ and $l \in \{1, 3, ...\}$, $\{y_{2nt+t+2ks+l}\}_{k\in\mathbb{N}}$ is a constant sequence eventually. Assume on the contrary that, for some $n \ge N$ and some $l \in \{1, 3, ...\}$, $\{y_{2nt+t+2ks+l}\}_{k\in\mathbb{N}}$ is not a constant sequence eventually. Then there exists a sequence of positive integers $k_1 < k_2 < \cdots$ such that, for any $r \in \mathbb{N}$, we have

$$B < y_{2nt+t+2k_rs+l} = \frac{A}{y_{2nt+t+2k_rs+l-s}}$$

$$< y_{2nt+t+2k_{r-1}s+l} = \frac{A}{y_{2nt+t+2k_{r-1}s+l-s}},$$
 (2.36)

which implies $y_{2nt+t+2k_rs+l-s} > y_{2nt+t+2k_{r-1}s+l-s}$ for any $r \in \mathbf{N}$. This is a contradiction. Take ps > N. Then $y_{2nst+t+l} = y_{2pst+t+2(nt-pt)s+l}$ is a constant sequence eventually for any $l \in \{1, 3, ...\}$. From the above we see that y_n is a periodic sequence with period 2st eventually.

In a similar fashion, we can show that if $x_{2nt+i} = A$ eventually for any $i \in \{0, 2, ...\}$ and $x_{2nt+l} \neq A$ eventually for some $l \in \{1, 3, ...\}$, then also statement (1(iii)) holds.

The second case follows from the previously proved one by interchanging letters. The proof is complete. $\hfill \Box$

Now we assume that A = B = 1. Then, for any $0 \le i \le 2t - 1$ and $n \in \mathbb{N}_0$, we have $1 \le x_{2(n+1)t+i} \le x_{2nt+i}$ eventually and $1 \le y_{2(n+1)t+i} \le y_{2nt+i}$ eventually.

Lemma 2.5 Let A = B = 1 and $s \ge t$. Then the following statements hold.

- (1) If $A_i = 1$, then $B_{t+i} = 1$. If $B_i = 1$, then $A_{t+i} = 1$.
- (2) If $x_N = 1$ for some $N \in \mathbb{N}$ and $A_{2nt+N+ks} = 1$ for any $k, n \in \mathbb{N}$, then $x_{2nt+N+ks} = y_{2nt+t+ks+N} = 1$ for any $k, n \in \mathbb{N}$. If $y_N = 1$ for some $N \in \mathbb{N}$ and $B_{2nt+N+ks} = 1$ for any $k, n \in \mathbb{N}$, then $y_{2nt+N+ks} = x_{2nt+t+ks+N} = 1$ for any $k, n \in \mathbb{N}$.
- (3) If s is even and gcd(s,t) = 1, then $1 \in \{x_n : n \in \{0,2,\ldots\}\} \cup \{y_{t+n} : n \in \{0,2,\ldots\}\}$ and $1 \in \{x_n : n \in \{1,3,\ldots\}\} \cup \{y_{t+n} : n \in \{1,3,\ldots\}\}.$
- (4) $1 \in \{x_n : n \in \mathbf{N}\} \cup \{y_n : n \in \mathbf{N}\}.$

Proof (1) Assume that $A_i = 1$. Assume on the contrary that $B_{t+i} > 1$. It follows from (1.1) that

$$y_{2nt+t+i} = \frac{x_{2nt+i}}{y_{2nt+t-s+i}}.$$
(2.37)

This implies

$$1 \le \lim_{n \to \infty} y_{2nt+t-s+i} = \frac{1}{B_{t+i}} < 1.$$
(2.38)

This is a contradiction. The second case follows from the previously proved one by interchanging letters.

(2) If $x_N = 1$ for some $N \in \mathbb{N}$, then $x_{2nt+N} = 1$ for any $n \in \mathbb{N}$. It follows from (1.1) that

$$y_{2nt+t+N} = \max\left\{1, \frac{x_{2nt+N}}{y_{2nt+t-s+N}}\right\}$$
$$= \max\left\{1, \frac{1}{y_{2nt+t-s+N}}\right\} = 1$$
(2.39)

and

$$y_{2nt+t+s+N} = \max\left\{1, \frac{x_{2nt+s+N}}{y_{2nt+t+N}}\right\}$$
$$= x_{2nt+s+N}$$
(2.40)

and

$$\begin{aligned} x_{2(n+1)t+N+s} &= \max\left\{1, \frac{y_{2nt+t+s+N}}{x_{2(n+1)t+N}}\right\} \\ &= \max\{1, y_{2nt+t+s+N}\} \\ &= y_{2nt+t+s+N} \\ &= x_{2nt+N+s}. \end{aligned}$$
(2.41)

Thus $x_{2nt+N+s} = y_{2nt+t+s+N} = 1$ for any $n \in \mathbb{N}$. In a similar fashion, we can show that $x_{2nt+N+ks} = y_{2nt+t+ks+N} = 1$ for any $k, n \in \mathbb{N}$. The second case follows from the previously proved one by interchanging letters.

(3) If *s* is even and gcd(s, t) = 1, then *t* is odd. Assume on the contrary that $1 \notin \{x_n : n \in \{0, 2, ...\}\} \cup \{y_{t+n} : n \in \{0, 2, ...\}\}$. Then it follows from (1.1) that, for any $n \in \mathbb{N}$,

$$y_{2nt+t} = \max\left\{1, \frac{x_{2nt}}{y_{2nt+t-s}}\right\} = \frac{x_{2nt}}{y_{2nt+t-s}} > 1$$
(2.42)

and

$$x_{2nt+2t-s} = \max\left\{1, \frac{y_{2nt+t-s}}{x_{2nt+2t-2s}}\right\}$$
$$= \frac{y_{2nt+t-s}}{x_{2nt+2t-2s}} > 1.$$
(2.43)

Thus

 $x_{2nt} > x_{2nt-2(s-t)}$ $> x_{2nt-4(s-t)}$ \dots $> x_{2t(n-t+s)}.$ (2.44)

This is a contradiction.

(4) Case (4) is treated similarly to case (3). The proof is complete.

Theorem 2.3 Let A = B = 1 and $s \ge t$. Then one of the following statements holds.

- (1) $x_n = 1$ eventually and y_n is a periodic sequence with period 2s eventually.
- (2) $y_n = 1$ eventually and x_n is a periodic sequence with period 2s eventually.
- (3) x_n , y_n are periodic sequences with period 2t eventually.

Proof If $x_n = 1$ (or $y_n = 1$) eventually, then by Lemma 2.1 we see that y_n (or x_n) is a periodic sequence with period 2*s* eventually. Now we assume that $x_n \neq 1$ eventually. Then we have $A_i > 1$ for some $0 \le i \le 2t - 1$ or $\lim_{n \to \infty} x_n = 1$.

If *s* is odd, then gcd(2t, s) = 1. Thus, for every $j \in \{0, 1, 2, ..., 2t - 1\}$, there exist some $1 \le i_j \le 2t$ and integer λ_j such that $i_j s = \lambda_j 2t + j$. By Lemma 2.3 and Lemma 2.5 we see that, for any $k \in \mathbb{N}$, $x_{2nt+ks+i}$ and $y_{2nt-t+ks+i}$ are constant sequences eventually, or for some $N \in \mathbb{N}$, $x_{2nt+N+ks} = y_{2nt+t+ks+N} = 1$ for any $k, n \in \mathbb{N}$, or for some $N \in \mathbb{N}$, $y_{2nt+N+ks} = x_{2nt+t+ks+N} = 1$ for any $0 \le r \le 2t - 1$, x_{2nt+r} and y_{2nt+r} are constant sequences eventually, which implies that x_n, y_n are periodic sequences with period 2t eventually.

In the following, we assume that *s* is even with s = 2s', then gcd(t, s') = 1 and *t* is odd. Thus, for every $j \in \{0, 1, 2, ..., t - 1\}$, there exist some $1 \le i_j \le t$ and integer λ_j such that $i_j s = \lambda_j 2t + 2j$.

If $A_i > 1$ for some $i \in \{0, 2, ...\}$, then by Lemma 2.3 we see that, for any $k \in \mathbb{N}$, $x_{2nt+ks+i}$ and $y_{2nt-t+ks+i}$ are constant sequences eventually. If $A_i = 1$ for any $i \in \{0, 2, ...\}$, then by Lemma 2.5 we have $B_{t+i} = 1$ for any $i \in \{0, 2, ...\}$ and $x_{2nt+i+ks} = y_{2nt-t+ks+i} = 1$ for any $k \in \mathbb{N}$ eventually. In a similar fashion, also we can show that, for any $i \in \{1, 3, ...\}$, $x_{2nt+ks+i}$ and $y_{2nt-t+ks+i}$ are constant sequences eventually for any $k \in \mathbb{N}$, or $x_{2nt+i+ks} = y_{2nt-t+ks+i} = 1$ for any $k \in \mathbb{N}$ eventually for any $i \in \{1, 3, ...\}$ and $k \in \mathbb{N}$. Thus, for any $0 \le r \le 2t - 1$, x_{2nt+r} and y_{2nt+r} are constant sequences eventually. This implies that x_n , y_n are periodic sequences with period 2t eventually.

Using the previously proved one by interchanging letters, also we can show that if $y_n \neq 1$ eventually, then x_n , y_n are periodic sequences with period 2t eventually. The proof is complete.

In Example 3.1 of [37], we showed that the equation

$$x_n = \frac{x_{n-t}}{x_{n-s}} (t > s)$$
(2.45)

has a positive solution z_n ($n \ge -t$) with $1 < z_{n+1} < z_n$ for any $n \ge -t$ and $\lim_{n \to \infty} z_n = 1$. From Example 3.1 of [37], we obtain the following theorem.

Theorem 2.4 Let $A \le 1$ and $B \le 1$ and s < t. Assume z_n $(n \ge -t)$ is a positive solution of (2.45) with $1 < z_{n+1} < z_n$ for any $n \ge -t$ and $\lim_{n\to\infty} z_n = 1$. Then equation (1.1) have a solution (x_n, y_n) with $1 < x_{n+1} = y_{n+1} = z_{n+1} < x_n = y_n = z_n$ for any $n \ge -t$ and $\lim_{n\to\infty} x_n = \lim_{n\to\infty} y_n = 1$.

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Competing interests

The authors declare that they have no competing interests.

Authors' contributions

All authors participated in every phase of research conducted for this paper. All authors read and approved the final manuscript.

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